

NEIL H. E. WESTE DAVID MONEY HARRIS

FOURTH EDITION

CMOS

Lecture 13: Datapath Functional Units

Outline

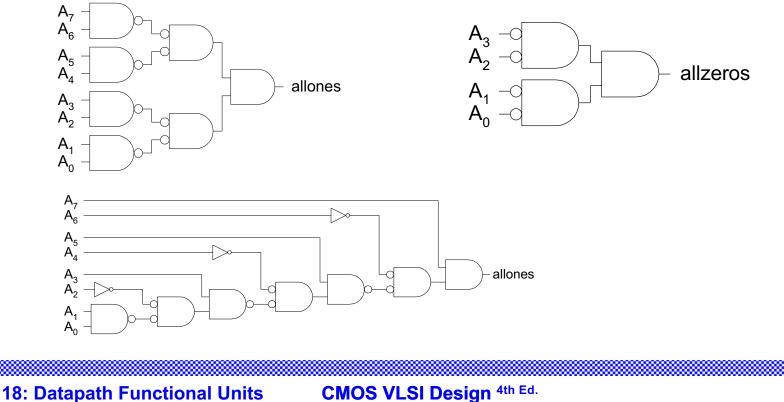
- Comparators
- Shifters
- Multi-input Adders
- Multipliers

Comparators

- **D** 0's detector: A = 00...000
- □ 1's detector: A = 11...111
- $\Box Equality comparator: A = B$
- □ Magnitude comparator: A < B

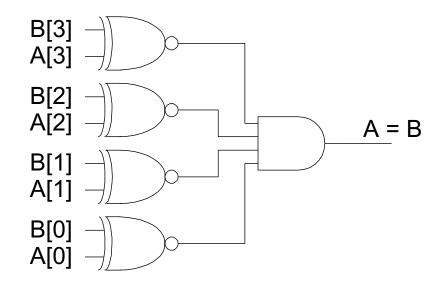
1's & 0's Detectors

- □ 1's detector: N-input AND gate
- □ 0's detector: NOTs + 1's detector (N-input NOR)



Equality Comparator

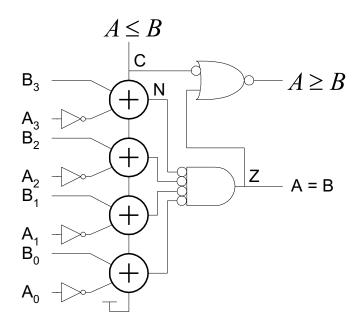
- □ Check if each bit is equal (XNOR, aka equality gate)
- □ 1's detect on bitwise equality



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Magnitude Comparator

- Compute B A and look at sign
- □ B A = B + ~A + 1
- For unsigned numbers, carry out is sign bit



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Signed vs. Unsigned

□ For signed numbers, comparison is harder

- C: carry out
- Z: zero (all bits of A B are 0)
- N: negative (MSB of result)
- V: overflow (inputs had different signs, output sign \neq B)
- S: N xor V (sign of result)

Relation	Unsigned Comparison	Signed Comparison
A = B	Ζ	Ζ
$A \neq B$	Z	Z
A < B	$C \cdot \overline{Z}$	$\overline{S} \cdot \overline{Z}$
A > B	\overline{C}	S
$A \leq B$	C	\overline{S}
$A \ge B$	$\overline{C} + Z$	S + Z

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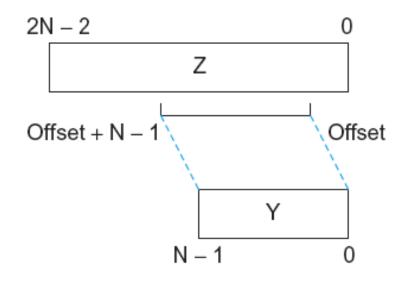
Shifters

□ Logical Shift:

- Shifts number left or right and fills with 0's
 - 1011 LSR 1 = 1011 LSL1 =
- □ Arithmetic Shift:
 - Shifts number left or right. Rt shift sign extends
 - 1011 ASR1 = 1011 ASL1 =
- Rotate:
 - Shifts number left or right and fills with lost bits
 - 1011 ROR1 = 1011 ROL1 =

Funnel Shifter

- A funnel shifter can do all six types of shifts
- Selects N-bit field Y from 2N–1-bit input
 - Shift by k bits $(0 \le k \le N)$
 - Logically involves N N:1 multiplexers



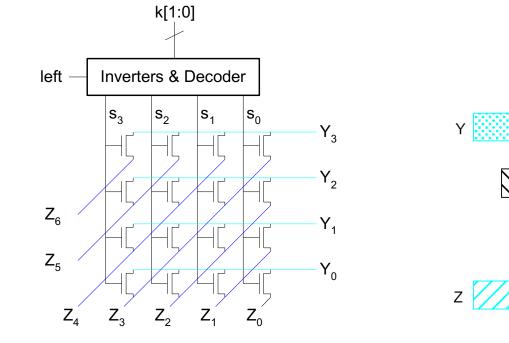
Funnel Source Generator

Shift Type	$Z_{2N-2:N}$	<i>Z_{N - 1}</i>	Z _{N-2:0}	Offset
Rotate Right	$A_{N-2:0}$	A_{N-1}	$A_{N-2:0}$	k
Logical Right	0	A_{N-1}	$A_{N-2:0}$	k
Arithmetic Right	sign	A_{N-1}	$A_{N-2:0}$	k
Rotate Left	A _{N-1:1}	A_0	$A_{N-1:1}$	k
Logical/Arithmetic Left	A _{N-1:1}	A_0	0	k

Array Funnel Shifter

N N-input multiplexers

- Use 1-of-N hot select signals for shift amount
- nMOS pass transistor design (V_t drops!)



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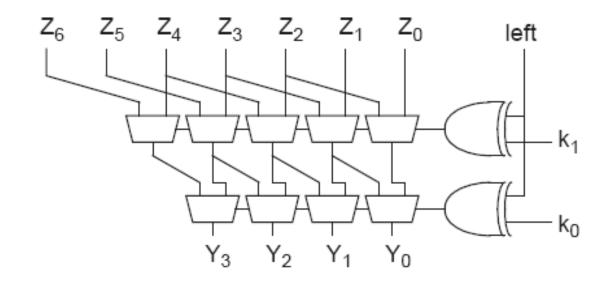
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Logarithmic Funnel Shifter

□ Log N stages of 2-input muxes

No select decoding needed

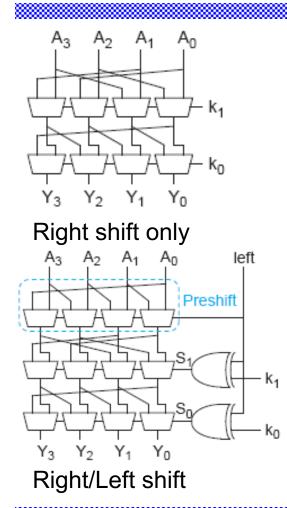


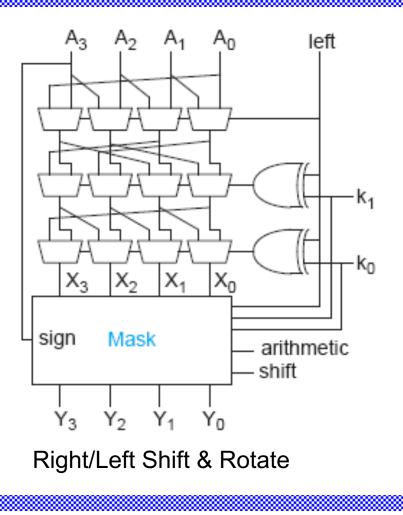
32-bit Logarithmic Funnel Wider multiplexers reduce delay and power Operands > 32 bits introduce datapath irregularity 32 Source 31 0 arith Generator Bits Source Generator shif 0-31left Z. 63 Source 31 Generator Bits 32-62 Mux 4:1 Decod 39 k 38 0 mux4 5 8:1 D 8 31 mux8 32 **18: Datapath Functional Units** CMOS VLSI Design 4th Ed. 13

Barrel Shifter

- Barrel shifters perform right rotations using wraparound wires.
- □ Left rotations are right rotations by N k = k + 1 bits.
- □ Shifts are rotations with the end bits masked off.

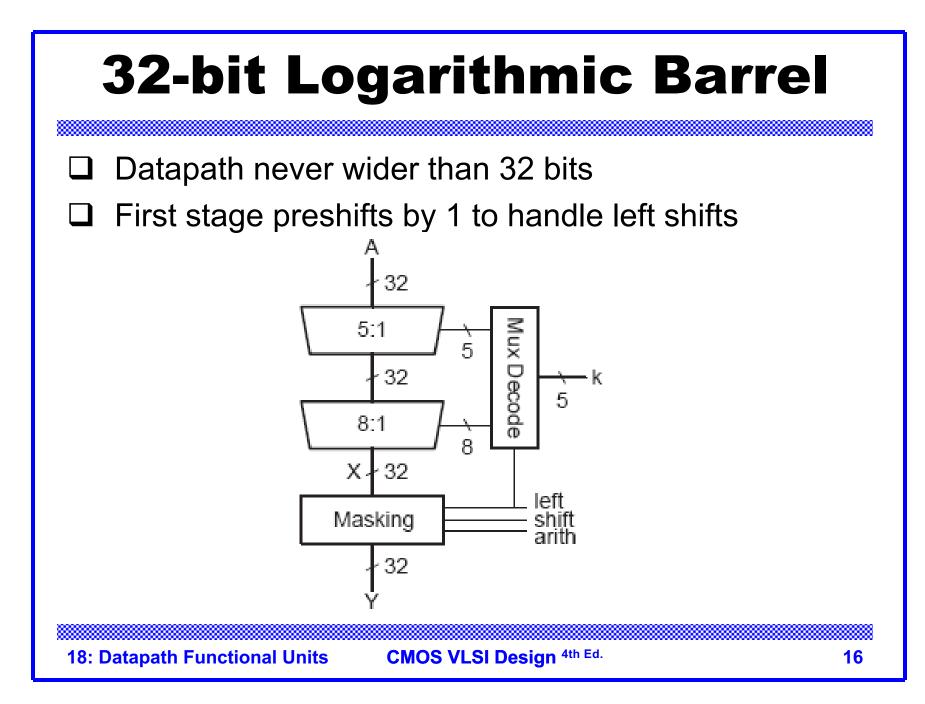
Logarithmic Barrel Shifter





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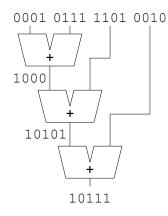
Multi-input Adders

□ Suppose we want to add k N-bit words

- Ex: 0001 + 0111 + 1101 + 0010 = 10111

□ Straightforward solution: k-1 N-input CPAs

Large and slow

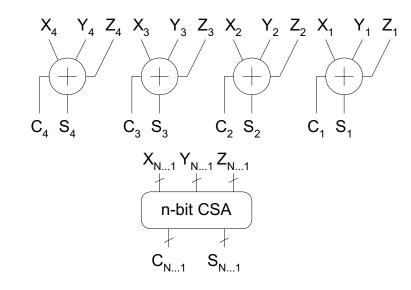


Carry Save Addition

□ A full adder sums 3 inputs and produces 2 outputs

- Carry output has twice weight of sum output

- □ N full adders in parallel are called *carry save adder*
 - Produce N sums and N carry outs



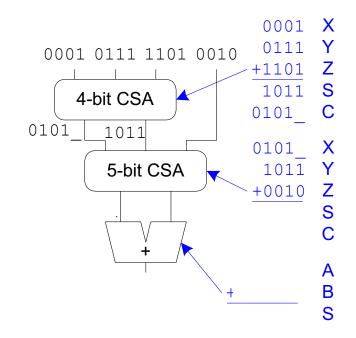
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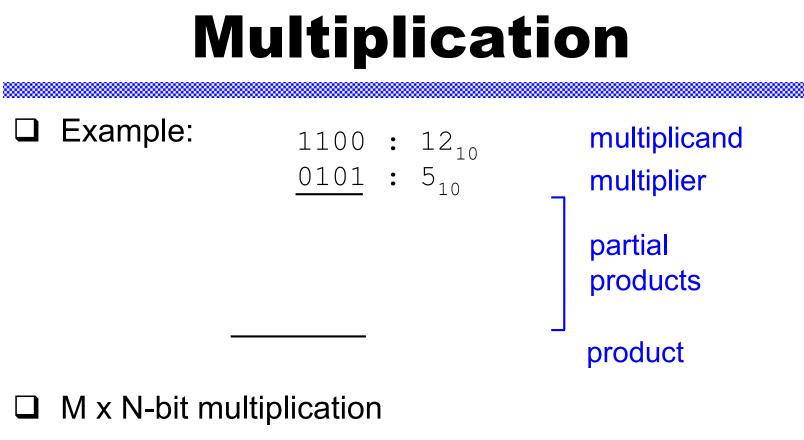
CSA Application

□ Use k-2 stages of CSAs

- Keep result in carry-save redundant form
- Final CPA computes actual result



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- Produce N M-bit partial products
- Sum these to produce M+N-bit product

General Form

Multiplicand: $Y = (y_{M-1}, y_{M-2}, ..., y_1, y_0)$ Multiplier: $X = (x_{N-1}, x_{N-2}, ..., x_1, x_0)$

D Product:
$$P = \left(\sum_{j=0}^{M-1} y_j 2^j\right) \left(\sum_{i=0}^{N-1} x_i 2^i\right) = \sum_{i=0}^{N-1} \sum_{j=0}^{M-1} x_i y_j 2^{i+j}$$

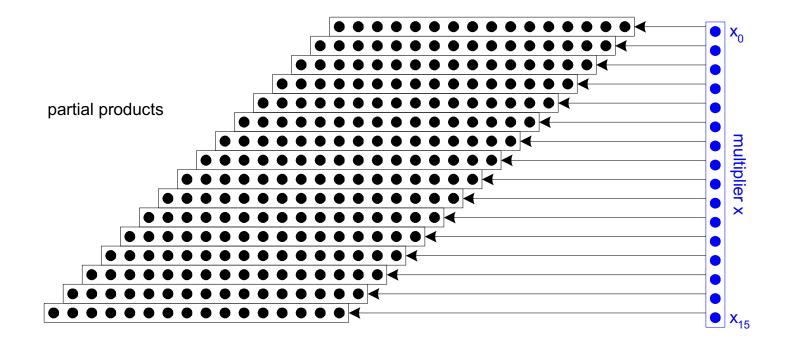
multiplicand У₅ $y_4 \quad y_3 \quad y_2 \quad y_1$ y₀ multiplier X₄ X₃ Х₅ х₂ **x**₁ X₀ $x_0y_5 x_0y_4 x_0y_3 x_0y_2 x_0y_1$ x₀y₀ $x_1y_5 \quad x_1y_4 \quad x_1y_3 \quad x_1y_2 \quad x_1y_1 \quad x_1y_0$ partial $x_2y_5 x_2y_4 x_2y_3 x_2y_2 x_2y_1 x_2y_0$ products $x_{3}y_{5}$ $x_{3}y_{4}$ $x_{3}y_{3}$ $x_{3}y_{2}$ $x_{3}y_{1}$ $x_{3}y_{0}$ $\mathbf{x}_4 \mathbf{y}_5 \quad \mathbf{x}_4 \mathbf{y}_4$ $x_4y_3 \quad x_4y_2 \quad x_4y_1 \quad x_4y_0$ x_5y_3 x₅y₂ $X_5 y_5$ x₅y₄ $\mathbf{x}_5 \mathbf{y}_1 \quad \mathbf{x}_5 \mathbf{y}_0$ product **p**₁₀ p₈ р₇ p_5 p_4 p_3 \mathbf{p}_2 p_1 \mathbf{p}_0 p₁₁ p₉ p_6

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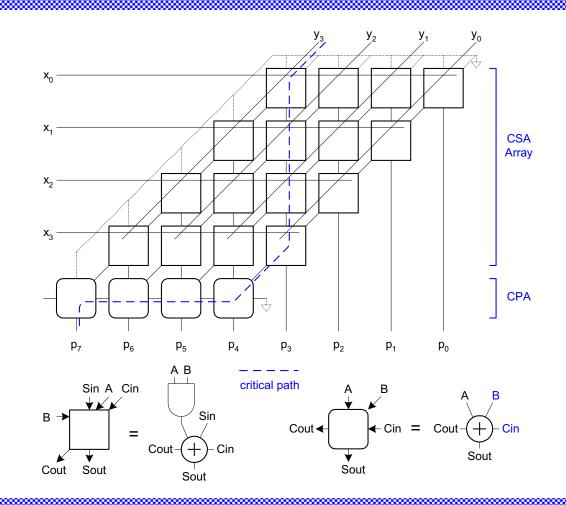
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Dot Diagram

Each dot represents a bit



Array Multiplier

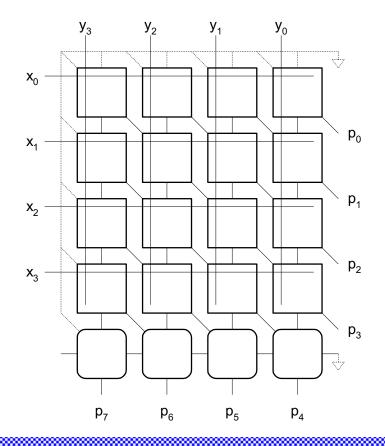


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Rectangular Array

□ Squash array to fit rectangular floorplan



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Fewer Partial Products

- Array multiplier requires N partial products
- If we looked at groups of r bits, we could form N/r partial products.
 - Faster and smaller?
 - Called radix-2^r encoding
- □ Ex: r = 2: look at pairs of bits
 - Form partial products of 0, Y, 2Y, 3Y
 - − First three are easy, but 3Y requires adder ⊗

Booth Encoding

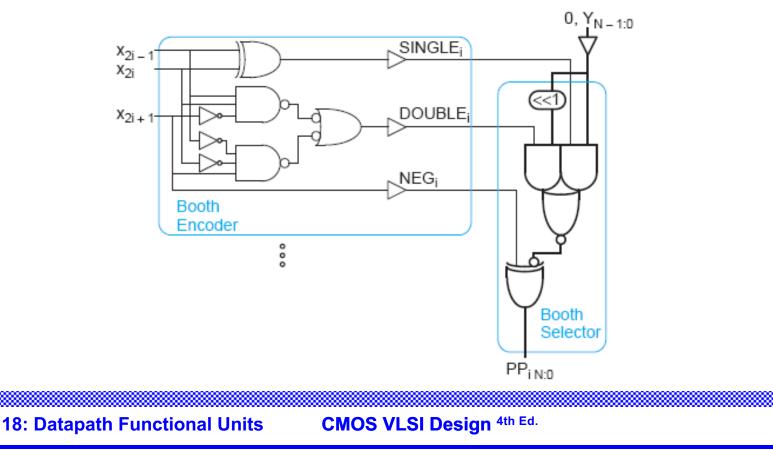
- Instead of 3Y, try –Y, then increment next partial product to add 4Y
- □ Similarly, for 2Y, try –2Y + 4Y in next partial product

Inputs		Partial Product	Booth Selects			
x _{2i+1}	x_{2i}	^x 2i−1	PP_i	$SINGLE_i$	DOUBLE_i	NEG_i
0	0	0	0	0	0	0
0	0	1	Y	1	0	0
0	1	0				
0	1	1				
1	0	0				
1	0	1				
1	1	0				
1	1	1				

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Booth Hardware

- □ Booth encoder generates control lines for each PP
 - Booth selectors choose PP bits

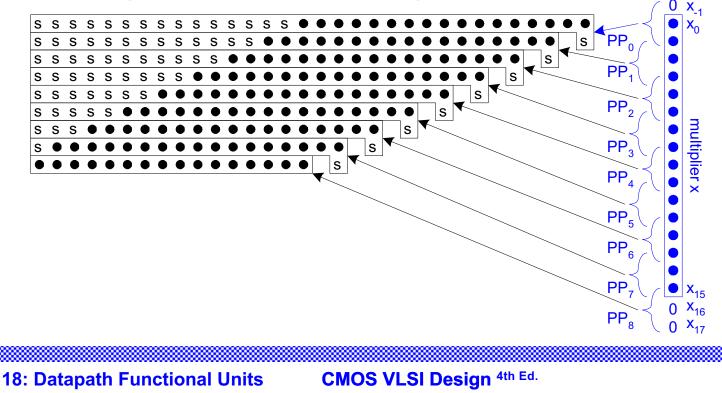


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Sign Extension

Partial products can be negative

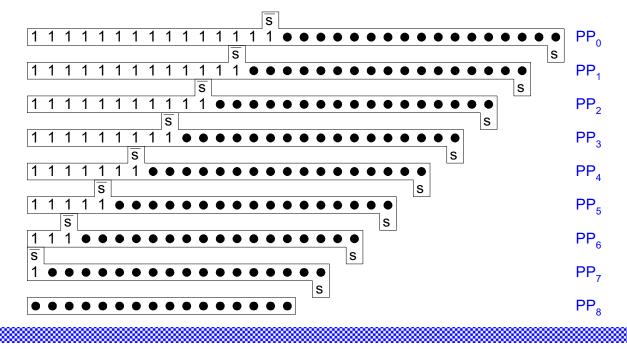
- Require sign extension, which is cumbersome
- High fanout on most significant bit



Simplified Sign Ext.

□ Sign bits are either all 0's or all 1's

- Note that all 0's is all 1's + 1 in proper column
- Use this to reduce loading on MSB

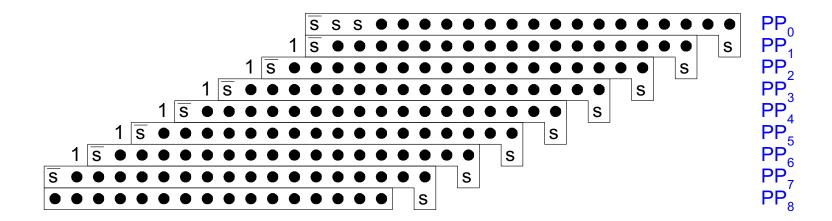


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Even Simpler Sign Ext.

□ No need to add all the 1's in hardware

- Precompute the answer!



Advanced Multiplication

- □ Signed vs. unsigned inputs
- Higher radix Booth encoding
- Array vs. tree CSA networks